

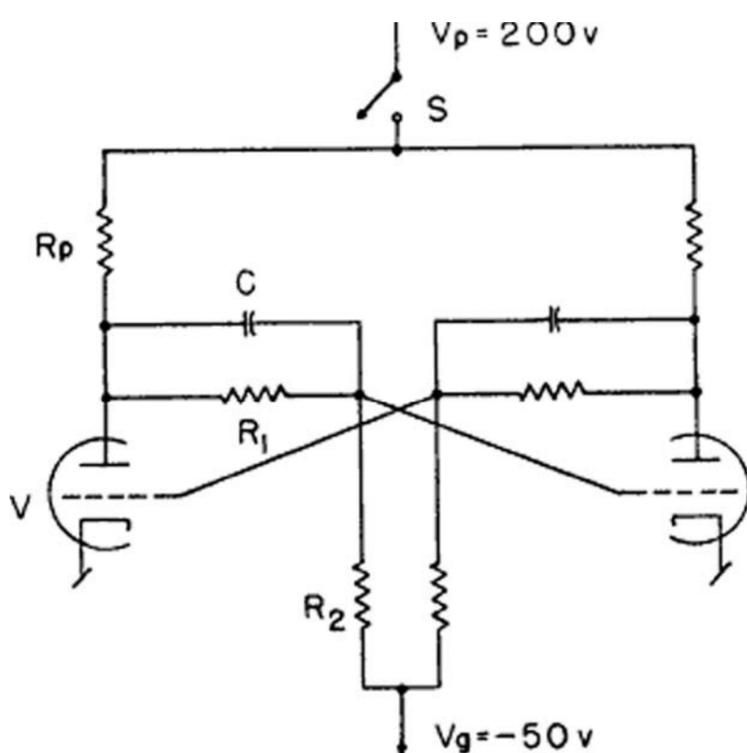
### Flip-flop as Generator of Random Binary Digits

The aim of the present note is to show that a well known electronic element of digital computers, **the flip-flop**, may be used for generating a series of random binary digits with equal probabilities.

Let us consider a **flip-flop** as shown on **fig. 1** and let **A** and **B** denote two possible stable states of the flip-flop. If we switch on the contact **S**, the flip-flop will be randomly set in one of its states **A** or **B**. We may obtain by the aid of the flip-flop a sequence of  $2k$  random elements  $X_1, X_2, \dots, X_{2k}$ , (abbreviated  $\{X_{2k}\}$ ), where

$$X_j = \begin{cases} \mathbf{A}, & \text{if } j\text{-th switching on the contact } S, \text{ set flip-flop in state } \mathbf{A} \\ \mathbf{B}, & \text{if } j\text{-th switching on the contact } S, \text{ set flip-flop in state } \mathbf{B} \end{cases}$$

and  $1 \leq j \leq 2k$ .



- $R_p = 10 \text{ K}\Omega$
- $R_1 = 160 \text{ K}\Omega$
- $R_2 = 50 \text{ K}\Omega$
- $C = 50 \mu\mu\text{fd}$
- $V = \frac{1}{2} \text{ 6SN7}$

FIG. 1.

In this way we may obtain a finite random series of **A** and **B** which are statistically independent. One series produced by the aid of a flip-flop is given below:

**AABAABBABBBABBBABBAABABABBABAABABABBAA**  
**BABABBBBBBBBBBABBBAABAABBB.**

Let  $\{Y_k\}$  be the sequence of  $k$  pairs of elements of  $\{X_{2k}\}$  such that  $Y_i = X_{2i-1}, X_{2i}$ , where  $1 \leq i \leq k$ . Omitting in  $\{Y_k\}$  all elements of the form **AA** and **BB** we obtain a third sequence whose elements are the pairs **AB** and **BA** only, denoted in the following by 0 and 1 respectively.

Let  $p_j(\mathbf{A})$  and  $p_j(\mathbf{B})$  denote probabilities that  $j$ -th switching on of contact **S** set flip-flop in state **A** or **B** respectively and suppose that  $p_j(\mathbf{A})$  and  $p_j(\mathbf{B})$  are asymmetric, say  $p_j(\mathbf{A}) > p_j(\mathbf{B})$ . Supposing that the flip-flop does not change its properties during two successive switchings, we may write

$$(1) \quad p_{2i-1}(\mathbf{A}) = p_{2i}(\mathbf{A})$$

$$(2) \quad p_{2i-1}(\mathbf{B}) = p_{2i}(\mathbf{B}).$$

From 1 and 2 we have

$$(3) \quad p_{2i-1}(\mathbf{A}) \cdot p_{2i}(\mathbf{B}) = p_{2i-1}(\mathbf{B}) \cdot p_{2i}(\mathbf{A}).$$

Because

$$(4) \quad p_{2i-1}(\mathbf{A}) \cdot p_{2i}(\mathbf{B}) = p_i(0)$$

and

$$(5) \quad p_{2i-1}(\mathbf{B}) \cdot p_{2i}(\mathbf{A}) = p_i(1),$$

therefore

$$(6) \quad p_i(0) = p_i(1)$$

where  $p_i(0)$  and  $p_i(1)$  are probabilities of zeros and ones in the  $i$ -th place of the sequence  $\{Y_k\}$ .

The procedure above described may be used for production of binary random numbers by automatic digital computers. In this case the manual switch **S** must be replaced by an electronic switch, of course.

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